

NEW HARDNESS RESULTS IN RAINBOW CONNECTIVITY

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ABSTRACT. A path in an edge colored graph is said to be a rainbow path if no two edges on the path have the same color. An edge colored graph is (strongly) rainbow connected if there exists a (geodesic) rainbow path between every pair of vertices. The (strong) rainbow connectivity of a graph G , denoted by $(src(G), rc(G))$, respectively) $rc(G)$ is the smallest number of colors required to edge color the graph such that the graph is (strong) rainbow connected. It is known that for *even* k to decide whether the rainbow connectivity of a graph is at most k or not is NP-hard. It was conjectured that for all k , to decide whether $rc(G) \leq k$ is NP-hard. In this paper we prove this conjecture. We also show that it is NP-hard to decide whether $src(G) \leq k$ or not even when G is a bipartite graph.

1. INTRODUCTION

This paper deals with the notion of *rainbow coloring* and *strong rainbow coloring* of a graph. Consider an edge coloring (not necessarily proper) of a graph $G = (V, E)$. We say that there exists a *rainbow path* between a pair of vertices, if no two edges on the path have the same color. A coloring of the edges in graph G is called a rainbow coloring if between every pair of vertices in G there exists a rainbow path. An edge coloring is a *strong* rainbow coloring if between every pair of vertices, one of its geodesic *i.e.*, shortest paths is a rainbow path. The minimum number of colors required to rainbow color a graph G is called the rainbow connection number denoted by $rc(G)$. Similarly, the minimum number of colors required to strongly rainbow color a graph G is called the strong rainbow connection number, denoted by $src(G)$. The notions of rainbow coloring and strong rainbow coloring were introduced recently by Chartrand et al. [5] as a means of strengthening the connectivity. Subsequent to this paper, the problem has received attention by several people and the complexity as well as upper bounds for the rainbow connection number have been studied.

In this paper we study the complexity of computing $rc(G)$ and $src(G)$. We prove the following results:

- (1) For every $k \geq 3$, deciding whether $src(G) \leq k$, is NP-hard even when G is bipartite.
- (2) Deciding rainbow connection number of a graph is at most 3 is NP-hard even when the graph G is bipartite.
- (3) For every $k \geq 3$, deciding whether $rc(G) \leq k$ is NP-hard.

We note that Chakraborty et al. [3] proved that for all even k , deciding whether $rc(G) \leq k$ is NP-hard. For proving the result they introduced the problem of *subset* rainbow connectivity where in addition to the graph $G = (V, E)$ we are given a set P containing pairs of vertices. The goal is to answer whether there exists an edge coloring of G with k colors such that every pair in P has a rainbow path. We also use the subset rainbow connectivity problem and analogously define the subset strong rainbow connectivity problem to prove our hardness results.

1.1. Related work. The concepts of rainbow connectivity and strong rainbow connectivity were first introduced by Chartrand et al. in [5]. There they computed the rainbow connection number and strong rainbow connection number for several graphs including the complete bipartite graph and multipartite graphs. Subsequent to this upper bounds for the rainbow connectivity as a function of minimum degree and the number of vertices of the graph were explored in [2],[8],[4]. Graphs of diameter 2 were studied in [9] and it was shown that $rc(G)$ is upper bounded by $k + 2$ where k is the number of bridges in the graph. Some upper bounds were shown in [4] for special graphs like interval graphs, AT-free graphs. The threshold function for random graph to have $rc(G) = 2$ was studied in [2]. In [8], the concept of rainbow vertex connection number was studied. In [1], Basavaraju et al. gave a constructive argument to show that any graph G can be colored with $r(r + 2)$ colors in polynomial time where r is the radius of the graph. In [3], it was shown that the following problem is NP-hard: Given a graph G and an even number $k(> 0)$, is $rc(G) \leq k$? It was conjectured that it is NP-hard to determine whether $rc(G) \leq k$ for all $k > 0$. Similar complexity results were shown for the rainbow vertex connection number in [6].

1.2. Organisation of the paper. We prove complexity results related to strong rainbow connectivity in Section 2. In Section 3, we give the complexity results related to the rainbow connectivity.

2. STRONG RAINBOW CONNECTIVITY

In this section we prove the hardness result for the following problem: given a graph G and an integer $k \geq 3$, decide whether $\text{src}(G) \leq k$. In order to show the hardness of this problem, we will first consider an intermediate problem called the k -subset strong rainbow connectivity problem which is the decision version of the subset strong rainbow connectivity problem. The input to the k -subset strong rainbow connectivity problem is a graph G along with a set of pairs $P = \{(u, v) : (u, v) \subseteq V \times V\}$ and an integer k . Our goal is to answer whether there exists an edge coloring of G with at most k colors such that every pair $(u, v) \in P$ has a geodesic rainbow path.

Our overall plan is to prove that k -subset strong rainbow connectivity is NP-hard by showing a reduction from the vertex coloring problem. We then establish the polynomial time equivalence of the k -subset strong rainbow connectivity problem and the problem of determining whether $\text{src}(G) \leq k$ for a graph G .

2.1. k -subset strong rainbow connectivity. Let $G = (V, E)$ be an instance of the k -vertex coloring problem. We say that G can be vertex colored using k colors if there exists an assignment of at most k colors to the vertices of G such that no pair of adjacent vertices are colored using the same color. This problem is NP-hard for $k \geq 3$. Given an instance $G = (V, E)$ of the k -vertex coloring problem, we construct an instance $\langle G' = (V', E'), P \rangle$ of the k -subset strong rainbow connectivity problem.

The graph G' that we construct is a star, with one leaf vertex corresponding to every vertex $v \in V$ and an additional central vertex a . The set of pairs P captures the edges in E , that is, for every edge $(u, v) \in E$ we have a pair (u, v) in the set P . The goal is to color the edges of G' using at most k colors such that every pair in the set P has a geodesic rainbow path. More formally, we define each of the parameters $\langle G' = (V', E'), P \rangle$ of the k -subset strong rainbow connectivity problem below:

$$\begin{aligned} V' &= \{a\} \cup V; & E' &= \{(a, v) : v \in V\} \\ P &= \{(u, v) : (u, v) \in E\}; \end{aligned}$$

We now prove the following lemma which establishes the hardness of the k -subset strong rainbow connectivity problem.

Lemma 2.1. *The graph $G = (V, E)$ is vertex colorable using $k(\geq 3)$ colors iff the graph $G' = (V', E')$ can be colored using k colors such*

that for every pair $(u, v) \in P$ there is a geodesic rainbow path between u and v .

Proof. Assume that G can be vertex colored using k colors; we now show an assignment of colors to the edges of the graph G' . Let c be the color assigned to a vertex $v \in V$; we assign the color c to the edge $(a, v) \in E'$. Now consider any pair $(u, v) \in P$. Recall that $(u, v) \in P$ because there exists an edge $(u, v) \in E$. Since the coloring was a proper vertex coloring, the edges (a, u) and (a, v) in G' are assigned different colors by our coloring. Thus, the path $\{(u, a), (a, v)\}$ is a rainbow path; further since that is the only path between u and v it is also a strong rainbow path.

To prove the other direction, assume that there exists an edge coloring of G' using k colors such that between every pair of vertices in P there is a rainbow path. It is easy to see that if we assign the color c of the edge $(a, v) \in E'$ to the vertex $v \in V$, we get a coloring that is a proper vertex coloring for G . \square \square

Recall the problem of subset rainbow connectivity where instead of asking for a geodesic rainbow path between every pair in P , we are content with any rainbow path. Note that our graph G' construed in the above reduction is a tree, in fact a star and hence between every pair of vertices there is exactly one path. Thus, all the above arguments apply for the k -subset rainbow connectivity problem as well. As a consequence we can conclude the following:

Lemma 2.2. *For every $k \geq 3$, k -subset strong rainbow connectivity and k -subset rainbow connectivity is NP-hard even when the input graph G is a star.*

From Lemma 2.1, it can be observed that if subset (strong) rainbow connectivity problem had a $f(n)$ -factor approximation algorithm then even the classical vertex coloring has a $f(n)$ -approximation algorithm. But we know from [7], that there is no $n^{1-\epsilon}$ factor approximation algorithm for vertex coloring unless $\text{NP}=\text{ZPP}$. Hence subset (strong) rainbow connectivity problem is inapproximable within a factor of $n^{1-\epsilon}$ unless $\text{NP}=\text{ZPP}$.

2.2. Hardness of k -strong rainbow connectivity. In this section we prove that the problem of deciding whether a graph G can be strongly rainbow colored using k colors is polynomial time equivalent to the k -subset strong rainbow connectivity problem *i.e.*, the existence of a polynomial time algorithm for one of them will imply the existence of polynomial time algorithm for the other. In particular we prove the lemma.

Lemma 2.3. *The following two problems are reducible to each other in polynomial time:*

- (1) *Given a graph $G = (V, E)$ and an integer k , decide whether the edges of G can be colored using k colors such that between every pair of vertices in G there is a geodesic rainbow path.*
- (2) *Given a graph $G = (V, E)$, an integer k and a set of pairs $P \subseteq V \times V$, decide whether the edges of G can be colored using k colors such that between every pair $(u, v) \in P$ there is a geodesic rainbow path.*

Proof. It suffices to prove that problem (2) reduces to problem (1). Let $\langle G = (V, E), P \rangle$ be an instance of the k -subset strong rainbow connectivity problem. Using Lemma 2.2, we know that k -subset strong rainbow connectivity is NP-hard even when G is a star as well as the pairs $(u, v) \in P$ are such that both u and v are leaf nodes of the star. We assume both these properties on the input $\langle G, P \rangle$ and use it crucially in our reduction. Let us denote the central vertex of the star G by a and the leaf vertices by $L = \{v_1, \dots, v_n\}$, that is, $V = \{a\} \cup L$. Using the graph G and the pairs P , we construct the new graph $G' = (V', E')$ as follows: for every leaf node $v_i \in L$, we introduce two new vertices u_i and u'_i . For every pair of leaf nodes $(v_i, v_j) \in (L \times L) \setminus P$, we introduce two new vertices $w_{i,j}$ and $w'_{i,j}$.

$$\begin{aligned} V' &= V \cup V_1 \cup V_2 \\ V_1 &= \{u_i : i \in \{1, \dots, n\}\} \cup \{w_{i,j} : (v_i, v_j) \in (L \times L) \setminus P\} \\ V_2 &= \{u'_i : i \in \{1, \dots, n\}\} \cup \{w'_{i,j} : (v_i, v_j) \in (L \times L) \setminus P\} \end{aligned}$$

The edge set E' is be defined as follows:

$$\begin{aligned} E' &= E \cup E_1 \cup E_2 \cup E_3 \\ E_1 &= \{(v_i, u_i) : v_i \in L, u_i \in V_1\} \cup \{(v_i, w_{i,j}), (v_j, w_{i,j}) : (v_i, v_j) \in (L \times L) \setminus P\} \\ E_2 &= \{(x, x') : x \in V_1, x' \in V_2\} \\ E_3 &= \{(a, x') : x' \in V_2\} \end{aligned}$$

We now prove that G' is strongly rainbow colorable using k colors iff there is an edge coloring of G (using k colors) such that every pair $(u, v) \in P$ has a strong rainbow path. To prove one direction, we first note that, for all pairs $(u, v) \in P$, there is a two length path (u, a, v) in G and this path is also present in G' . Further, this path is the only two length path in G' between u and v ; hence any strong rainbow coloring of G' using k colors must make this path a rainbow path. This implies that if G cannot be edge colored with k colors such that every pair

in P is strongly rainbow connected, the graph G' cannot be strongly rainbow colored using k colors.

To prove the other direction, assume that there is a coloring $\chi : E \rightarrow \{c_1, c_2, \dots, c_k\}$ of G such that all pairs in P are strongly rainbow connected. We extend this coloring to a strong rainbow coloring of G' as follows:

- We retain the color on the edges of G , i.e. $\chi'(e) = \chi(e) : e \in E$.
- For each edge $(v_i, u_i) \in E_1$, we set $\chi'(v_i, u_i) = c_3$.
- For each pair of edges $\{(v_i, w_{i,j}), (v_j, w_{i,j})\} \in E_1$, we set $\chi'(v_i, w_{i,j}) = c_1$, $\chi'(v_j, w_{i,j}) = c_2$ (Assume without loss of generality that $i < j$).
- The edges in E_2 form a complete bipartite graph between the vertices in V_1 and V_2 . To color these edges, we pick a perfect matching M of size $|V_1|$ and assign $\chi'(e) = c_1, \forall e \in E_2 \cap M$ and $\chi'(e) = c_2, \forall e \in E_2 \setminus M$.
- Finally, for each edge $(a, x') \in E_3$, we set $\chi'(a, x') = c_3$.

It is straightforward to verify that this coloring is indeed a strong rainbow coloring of the graph G' using no more than k colors. This finishes the proof. □ □

We note that the graph $G' = (V', E')$ constructed in the proof of Lemma 2.3 is in fact bipartite. The vertex set V' can be partitioned into two sets A and B , where $A = \{a\} \cup V_1$ and $B = L \cup V_2$ such that there are not edges between vertices in the same partition. Thus, using Lemma 2.2 and Lemma 2.3, we conclude the following theorem.

Theorem 2.4. *For every $k \geq 3$, deciding whether a given graph G can be strongly rainbow colored using k colors is NP-hard. Further, the hardness holds even when the graph G is bipartite.*

From the same construction, it can be proved that deciding whether a given graph G can be rainbow colored using at most 3 colors is NP-hard.

3. RAINBOW CONNECTION NUMBER

In this section we investigate the complexity of deciding whether the rainbow connection number of a graph is less than or equal to some natural number k . We prove the following problem is NP-hard: given a graph G and an integer k decide whether $rc(G) \leq k$. We recall from Lemma 2.2 that the k -subset rainbow connectivity problem is NP-hard. We reduce from the k -subset rainbow connectivity to our problem of deciding whether a given graph can be rainbow colored using at most k colors.

3.1. Hardness of k -rainbow connectivity. We first describe the outline of our reduction. Let $\langle G = (V, E), P \rangle$ be the input to the k -subset rainbow connectivity problem where $P \subseteq V \times V$. We construct a graph $G' = (V', E')$ such that G is a subgraph of G' and $rc(G') \leq k$ iff G is k -subset rainbow connected. To construct G' , we first construct a graph $H_k = (W_k, E_k)$ such that $V \subset W_k$ and $V' = W_k$. Corresponding to the set P of pairs of vertices we associate a set of pairs of vertices P_k with respect to the graph H_k (The set P_k is only a relabelling of pairs of vertices in the set P). Then, the edge disjoint union of the two graphs G and H_k will yield the graph G' . The graph G' is constructed such that it satisfies two properties:

- (1) There exists an edge coloring χ of H_k with k colors such that all pairs of vertices in G' except those in P *i.e.*, all pairs in $(V' \times V') \setminus P$ are rainbow connected. Thus, if G is k -subset rainbow connected then there exists k -edge coloring of G such that all pairs in P will be rainbow connected. From this we prove that G' can be rainbow colored using k colors if G is k -subset rainbow connected.
- (2) All paths of length k or less between any pair of vertices in P are contained entirely in G (as a subgraph of G') itself. This ensures that for a rainbow coloring of G' with k colors, any pair of vertices in P should have all its rainbow paths inside G itself. Hence, if G' can be rainbow colored with k colors then G can be edge colored with k colors such that it is subset rainbow connected with respect to P .

We construct the family of graphs H_k inductively. For the base cases $k = 2$ and $k = 3$ we give explicit constructions and prove the properties mentioned above are true for the base cases. We then show our inductive step and prove that the properties hold. Finally, we describe our graph G' and prove the correctness of the reduction.

3.1.1. Construction of H_2 : The construction of the graph H_2 is derived from the reduction of Chakraborty et al. [3] used to prove that the 2-subset rainbow connectivity problem is NP-hard. Let $H_2 = (W_2, E_2)$ where the vertex set W_2 is defined as follows:

$$\begin{aligned}
 W_2 &= W_2^{(1)} \cup W_2^{(2)} \\
 W_2^{(1)} &= \{v_{i,2} : i \in \{1, \dots, n\}\} \\
 W_2^{(2)} &= \{u_i : i \in \{1, \dots, n\}\} \cup \{w_{i,j} : (v_i, v_j) \in (V \times V) \setminus P\}
 \end{aligned}$$

The edge set E_2 is defined as:

$$\begin{aligned} E_2 &= E_2^{(1)} \cup E_2^{(2)} \cup E_2^{(3)} \\ E_2^{(1)} &= \{(v_{i,2}, u_i) : i \in \{1, \dots, n\}\} \\ E_2^{(2)} &= \{(v_{i,2}, w_{i,j}), (v_{j,2}, w_{i,j}) : (v_i, v_j) \in (V \times V) \setminus P\} \\ E_2^{(3)} &= \{(x, y) : x, y \in W_2^{(2)}\} \end{aligned}$$

The set of vertices in $W_2^{(1)}$ are referred to as base vertices of H_2 . Let $P_2 = \{(v_{i,2}, v_{j,2}) : (v_i, v_j) \in P\}$. The graph H_2 has the property that for all $(v_{i,2}, v_{j,2}) \in P_2$ there is no path of length ≤ 2 between $v_{i,2}$ and $v_{j,2}$. Also, if $(v_{i,2}, v_{j,2}) \notin P_2$ the shortest path between $v_{i,2}$ and $v_{j,2}$ is of length 2.

Let the edge coloring $\chi : E_2 \rightarrow \{c_1, c_2\}$ of H_2 be defined as follows.

- If $e \in E_2^{(1)}$ then $\chi(e) = c_2$.
- For each pair of edges $\{(v_{i,2}, w_{i,j}), (v_{j,2}, w_{i,j})\} \in E_2^{(2)}$, we set $\chi((v_{i,2}, w_{i,j})) = c_1$ and $\chi((v_{j,2}, w_{i,j})) = c_2$ assuming without loss of generality that $i < j$.
- If $e \in E_2^{(3)}$ then $\chi(e) = c_1$.

It can be verified that for every pair of vertices (x, y) in $(W_2 \times W_2) \setminus P_2$, there is a rainbow path between x and y .

3.1.2. *Construction of H_3* : We now describe the construction of the graph H_3 . Let $H_3 = (W_3, E_3)$ be a graph where the vertex set W_3 is defined as follows:

$$\begin{aligned} W_3 &= W_3^{(1)} \cup W_3^{(2)} \cup W_3^{(3)} \\ W_3^{(1)} &= \{v_{i,3} : i \in \{1, \dots, n\}\} \\ W_3^{(2)} &= \{u_i : v_i \in V\} \cup \{a_{i,j}, b_{i,j} : (v_i, v_j) \in (V \times V) \setminus P\} \\ W_3^{(3)} &= \{u'_i : v_i \in V\} \cup \{a'_{i,j}, b'_{i,j} : (v_i, v_j) \in (V \times V) \setminus P\} \end{aligned}$$

The edge set E_3 is defined as:

$$\begin{aligned} E_3 &= E_3^{(1)} \cup E_3^{(2)} \cup E_3^{(3)} \cup E_3^{(4)} \\ E_3^{(1)} &= \{(v_{i,3}, u_i) : v_i \in V\} \\ E_3^{(2)} &= \{(v_{i,3}, a_{i,j}), (v_{j,3}, b_{i,j}) : (v_i, v_j) \in (V \times V) \setminus P\} \\ E_3^{(3)} &= \{(x, x') : x \in W_3^{(2)}, x' \in W_3^{(3)}\} \\ E_3^{(4)} &= \{(a_{i,j}, b_{i,j}) : (v_i, v_j) \in (V \times V) \setminus P\} \end{aligned}$$

The set of vertices in $W_3^{(1)}$ are referred to as base vertices of H_3 . Let $P_3 = \{(v_{i,3}, v_{j,3}) : (v_i, v_j) \in P\}$. The graph H_3 has the property that for all $(v_{i,3}, v_{j,3}) \in P_3$ there is no path of length ≤ 3 between $v_{i,3}$ and $v_{j,3}$. Also, if $(v_{i,3}, v_{j,3}) \notin P_3$ then the shortest path between $v_{i,3}$ and $v_{j,3}$ is of length 3.

Define an edge coloring $\chi : E_3 \rightarrow \{c_1, c_2, c_3\}$ of H_3 as follows.

- For each edge $e = (v_{i,3}, u_i) \in E_3^{(1)}$, we set $\chi(e) = c_3$.
- For each pair of edges $\{(v_{i,3}, a_{i,j}), (v_{j,3}, b_{i,j})\} \in E_3^{(2)}$, we set $\chi((v_{i,3}, a_{i,j})) = c_1$, $\chi((v_{j,3}, b_{i,j})) = c_2$ assuming without loss of generality that $i < j$.
- The edges in $E_3^{(3)}$ form a complete bipartite graph between the vertices in $W_3^{(2)}$ and $W_3^{(3)}$. To color these edges, we pick a perfect matching M of size $|W_3^{(2)}|$ and assign $\chi(e) = c_1, \forall e \in E_3^{(3)} \cap M$ and $\chi(e) = c_2, \forall e \in E_3^{(3)} \setminus M$.
- For an edge $e \in E_3^{(4)}$, assign $\chi(e) = c_3$.

It can be verified that for every pair of vertices (x, y) in $(W_3 \times W_3) \setminus P$, there is a rainbow path between x and y .

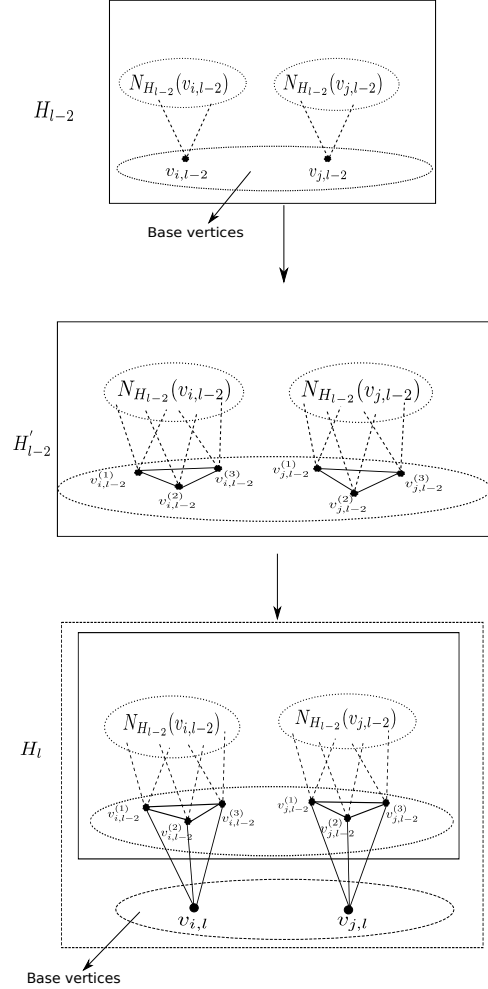
3.1.3. Inductive step: Assume we have constructed $H_{l-2} = (W_{l-2}, E_{l-2})$. From H_{l-2} we first construct an intermediate graph \mathcal{H}'_{l-2} from which we construct H_l . The graph \mathcal{H}'_{l-2} is constructed as follows.

Construction of \mathcal{H}'_{l-2} : Each base vertex $v_{i,l-2}$ in W_{l-2} is split into three vertices $v_{i,l-2}^{(1)}, v_{i,l-2}^{(2)}, v_{i,l-2}^{(3)}$ and edges are added between them *i.e.*, the vertices $v_{i,l-2}^{(1)}, v_{i,l-2}^{(2)}, v_{i,l-2}^{(3)}$ form a triangle. Any edge of the form $(w, v_{i,l-2})$ is replaced by three edges $(w, v_{i,l-2}^{(1)}), (w, v_{i,l-2}^{(2)}), (w, v_{i,l-2}^{(3)})$. Formally the graph $\mathcal{H}'_{l-2} = (\mathcal{W}'_{l-2}, \mathcal{E}'_{l-2})$ is defined as follows. The set of vertices is $\mathcal{W}'_{l-2} = \mathcal{W}_{l-2}^{(1)} \cup \mathcal{W}_{l-2}^{(2)}$ where,

$$\begin{aligned} \mathcal{W}_{l-2}^{(1)} &= \{v_{i,l-2}^{(1)}, v_{i,l-2}^{(2)}, v_{i,l-2}^{(3)} : i \in \{1, \dots, n\}\} \\ \mathcal{W}_{l-2}^{(2)} &= W_{l-2} \setminus \{v_{i,l-2} : i \in \{1, \dots, n\}\} \end{aligned}$$

and the edge set is $\mathcal{E}'_{l-2} = \mathcal{E}_{l-2}^{(1)} \cup \mathcal{E}_{l-2}^{(2)} \cup \mathcal{E}_{l-2}^{(3)}$ where,

$$\begin{aligned} \mathcal{E}_{l-2}^{(1)} &= \{(v_{i,l-2}^{(j_1)}, w) : (v_{i,l-2}, w) \in E_{l-2}, j_1 \in \{1, 2, 3\}\} \\ \mathcal{E}_{l-2}^{(2)} &= \{(v_{i,l-2}^{(j_1)}, v_{i,l-2}^{(j_2)}) : j_1, j_2 \in \{1, 2, 3\}\} \\ \mathcal{E}_{l-2}^{(3)} &= E_{l-2} \setminus \{(v_{i,l-2}, w) : w \in W_{l-2}\} \end{aligned}$$

FIGURE 1. Construction of H_l .

The graph $H_l = (W_l, E_l)$ where $W_l = \mathcal{W}'_{l-2} \cup V_l$ and $E_l = \mathcal{E}'_{l-2} \cup E'$ such that

$$\begin{aligned} V_l &= \{v_{1,l}, \dots, v_{n,l}\} \\ E' &= \{(v_{i,l-2}^{(1)}, v_{i,l}), (v_{i,l-2}^{(2)}, v_{i,l}), (v_{i,l-2}^{(3)}, v_{i,l}) : 1 \leq i \leq n\} \end{aligned}$$

The vertices in V_l are the base vertices of H_l . We now prove that the graph H_l satisfies the two properties stated at the beginning of the section. Before that, we define P_l to be the set $\{(v_{i_1,l}, v_{i_2,l}) : (v_{i_1}, v_{i_2}) \in P\}$ which is a set of pairs of vertices in H_l .

Lemma 3.1. *If $(v_{i_1,l}, v_{i_2,l}) \in P_l$ then there is no path between $v_{i_1,l}$ and $v_{i_2,l}$ of length less than l in G' .*

Proof. It was shown that the assertion was true for H_2 and H_3 . Assume that the assertion is true for H_{l-2} . Let $(v_{i_1}, v_{i_2}) \in P$. Then, $(v_{i_1, l-2}, v_{i_2, l-2}) \in P_{l-2}$ and hence by induction, every path from $v_{i_1, l-2}$ to $v_{i_2, l-2}$ in H_{l-2} is of length at most $l-2$. It can be seen that by the construction of \mathcal{H}'_{l-2} that we are not shortening the paths between any two vertices and all paths from $v_{i_1, l-2}^{(j_1)}$ to $v_{i_2, l-2}^{(j_2)}$ will still be of length at most $l-2$ for $j_1, j_2 \in \{1, 2, 3\}$. Consider the graph H_l . Since, $N_{H_l}(v_{i_1, l}) = \{v_{i_1, l-2}^{(1)}, v_{i_1, l-2}^{(2)}, v_{i_1, l-2}^{(3)}\}$ and $N_{H_l}(v_{i_2, l}) = \{v_{i_2, l-2}^{(1)}, v_{i_2, l-2}^{(2)}, v_{i_2, l-2}^{(3)}\}$ (here, $N_G(v)$ denotes the neighbourhood of vertex v in graph G), there cannot be path of length less than or equal to l between $v_{i_1, l}$ and $v_{i_2, l}$ in H_l . \square \square

Lemma 3.2. *If $(v_{i_1, l}, v_{i_2, l}) \notin P_l$ then the shortest path between $v_{i_1, l}$ and $v_{i_2, l}$ is of length l .*

Proof. It was shown that the assertion was true for H_2 and H_3 . Assume that the assertion is true for H_{l-2} . Let $v_{i_1}, v_{i_2} \notin P$. Then, $(v_{i_1, l-2}, v_{i_2, l-2}) \in P_{l-2}$ and hence by induction, the shortest path from $v_{i_1, l-2}$ to $v_{i_2, l-2}$ is of length $l-2$. The construction of \mathcal{H}'_{l-2} does not shorten any path between any two vertices and hence the shortest path from $v_{i_1, l-2}^{(j_1)}$ to $v_{i_2, l-2}^{(j_2)}$ will be of length exactly $l-2$ for $j_1, j_2 \in \{1, 2, 3\}$. Now consider the graph H_l . Since, $N_{H_l}(v_{i_1, l}) = \{v_{i_1, l-2}^{(1)}, v_{i_1, l-2}^{(2)}, v_{i_1, l-2}^{(3)}\}$ and $N_{H_l}(v_{i_2, l}) = \{v_{i_2, l-2}^{(1)}, v_{i_2, l-2}^{(2)}, v_{i_2, l-2}^{(3)}\}$, the shortest path $v_{i_1, l}$ and $v_{i_2, l}$ in H_l is of length l . \square \square

Lemma 3.3. *There exists an edge coloring of H_l with l colors such that all pairs in $(W_l \times W_l) \setminus P_l$ are rainbow connected.*

Proof. It was shown that the assertion is true for H_2, H_3 . Assume that the assertion is true for H_{l-2} . Let $\chi : E_{l-2} \rightarrow \{c_1, \dots, c_{l-2}\}$ be an edge coloring of H_{l-2} such that all pairs in $(W_{l-2} \times W_{l-2}) \setminus P_{l-2}$ are rainbow connected. Using χ , we define an edge coloring $\chi' : E_l \rightarrow \{c_1, \dots, c_l\}$ of H_l as follows.

- If $e \in \mathcal{E}_{l-2}^{(1)}$, then if $e = (v_{i, l-2}^{(1)}, w)$ or $e = (v_{i, l-2}^{(2)}, w)$ then $\chi'(e) = \chi((v_{i, l-2}, w))$ or else if $e = (v_{i, l-2}^{(3)}, w)$ then $\chi'(e) = c_{l-1}$.
- If $e \in \mathcal{E}_{l-2}^{(2)}$, then $\chi'(e) = c_l$.
- If $e \in \mathcal{E}_{l-2}^{(3)}$, then $\chi'(e) = \chi(e)$.
- If $e \in E'$, then if $e = (v_{i, l-2}^{(1)}, v_{i, l})$ then $\chi'(e) = c_{l-1}$ or else if $e = (v_{i, l-2}^{(2)}, v_{i, l})$ or $e = (v_{i, l}^{(3)}, v_{i, l})$ then $\chi'(e) = c_l$.

We now show that the l -edge coloring χ' makes H_l rainbow connected.

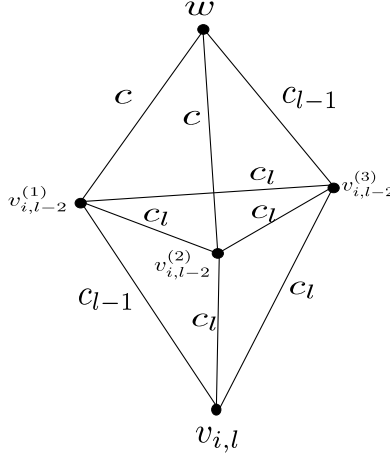


FIGURE 2. Edge coloring of H_l : c is the color originally used to color the edge $(v_{i,l-2}, w)$.

Claim. Let $v, w \in W_l$. If $(v, w) \notin P_l$, then v and w are rainbow connected in H_l .

Proof. We will consider the following cases.

Case (i) $v, w \in \mathcal{W}_{l-2}^{(2)}$: Since $\mathcal{W}_{l-2}^{(2)} \subset W_{l-2}$, there is a rainbow path \mathcal{P} from v to w in W_{l-2} . If \mathcal{P} consisted of edges only from $\mathcal{E}_{l-2}^{(3)}$ then these edges are also in E_{l-2} and by our coloring scheme, the rainbow path \mathcal{P} is also present in H_l . So assume that \mathcal{P} contains base vertices *i.e.*, let $\mathcal{P} = v \rightsquigarrow v_{i_1, l-2} \cdots \rightsquigarrow v_{i_s, l-2} \rightsquigarrow w$ where the distinct i_1, \dots, i_s belong to $\{1, \dots, n\}$. It can be observed that the path $\mathcal{P}' = v \rightsquigarrow v_{i_1, l-2}^{(1)} \cdots \rightsquigarrow v_{i_s, l-2}^{(1)} \rightsquigarrow w$ is rainbow path from v to w in W_{l-2} .

Case (ii) $v \in \mathcal{W}_{l-2}^{(1)}, w \in \mathcal{W}_{l-2}^{(2)}$: Let $v = v_{i, l-2}^{(j_1)}$ for $i \in \{1, \dots, n\}, j_1 \in \{1, 2, 3\}$. There is a rainbow path in H_{l-2} from $v_{i, l-2}$ to w . From an argument similar to the previous case, there exists a corresponding rainbow path from $v = v_{i, l-2}^{(1)}$ to w denoted by $v_{i, l-2}^{(1)} \rightsquigarrow^{R_{H_l}} w$ which does not use the colors c_{l-1} and c_l . Similarly there exists a rainbow path from $v_{i, l-2}^{(2)}$ to w . Also, the path $v_{i, l-2}^{(3)} - v_{i, l-2}^{(1)} \rightsquigarrow^{R_{H_l}} w$ from $v_{i, l-2}^{(3)}$ to w is a rainbow path.

Case (iii) $v, w \in \mathcal{W}_{l-2}^{(1)}$: Let $v = v_{i_1, l-2}^{(j_1)}, w = v_{i_2, l-2}^{(j_2)}$ for $j_1, j_2 \in \{1, 2, 3\}$. And let $w_1 \in N_{H_{l-2}}(v_{i_1, l-2}), w_2 \in N_{H_{l-2}}(v_{i_2, l-2})$. From Case (ii), there are rainbow paths from w_1 to $v_{j_1, l-2}^{(1)}$ (denoted by $w_1 \rightsquigarrow^{R_G} v_{j_1, l-2}^{(1)}$) and from w_2 to $v_{i_2, l-2}^{(1)}$ (denoted by $w_2 \rightsquigarrow^{R_G} v_{i_2, l-2}^{(1)}$). The following table

shows the rainbow paths for all possible cases of v and w :

	$v_{j,l-2}^{(1)}$	$v_{j,l-2}^{(2)}$	$v_{j,l-2}^{(3)}$
$v_{i,l-2}^{(1)}$	$v_{i,l-2}^{(1)} \overset{R_{H_l}}{\rightsquigarrow} w_2 - v_{j,l-2}^{(3)} - v_{j,l-2}^{(1)}$	$v_{i,l-2}^{(1)} \overset{R_{H_l}}{\rightsquigarrow} w_2 - v_{j,l-2}^{(3)} - v_{j,l-2}^{(2)}$	$v_{i,l-2}^{(1)} \overset{R_{H_l}}{\rightsquigarrow} w_2 - v_{j,l-2}^{(3)}$
$v_{i,l-2}^{(2)}$	—	$v_{i,l-2}^{(2)} \overset{R_{H_l}}{\rightsquigarrow} w_2 - v_{j,l-2}^{(3)} - v_{j,l-2}^{(2)}$	$v_{i,l-2}^{(2)} \overset{R_{H_l}}{\rightsquigarrow} w_2 - v_{j,l-2}^{(3)}$
$v_{i,l-2}^{(3)}$	—	—	$v_{i,l-2}^{(3)} - w_1 \overset{R_{H_l}}{\rightsquigarrow} v_{i,l-2}^{(1)} - v_{j,l-2}^{(3)}$

Case (iv) $v, w \in V_l$ and $(v, w) \notin P_l$: Let $v = v_{i,l}, w = v_{i,l}$. Since $(v_{i,l}, v_{j,l}) \notin P_l$, this means that $(v_{i,l-2}, v_{j,l-2}) \notin P_{l-2}$. Hence, there exists a rainbow path from $v_{i,l-2}$ to $v_{j,l-2}$ in H_{l-2} . Correspondingly in the graph \mathcal{H}'_{l-2} , there exists a rainbow path from $v_{i,l-2}^{(1)}$ to $v_{j,l-2}^{(2)}$ denoted by $v_{i,l-2}^{(1)} \rightsquigarrow v_{j,l-2}^{(2)}$. It can be observed that the edge colored path $v_{i,l-2} - v_{i,l-2}^{(1)} \rightsquigarrow v_{j,l-2}^{(2)} - v_{j,l-2}$ is a rainbow path.

Case (v) $v \in W_{l-2}, w \in V_l$: Let $w = v_{i,l} \in V_l$. There exists a rainbow path from $v_{i,l-2}$ to v in H_{l-2} . Correspondingly there exists a rainbow path from $v_{i,l-2}^{(1)}$ to v in H_l , denoted by $v_{i,l-2}^{(1)} \rightsquigarrow v$, which does not use colors c_{l-1} and c_l . Hence, the path $v_{i,l} - v_{i,l-2}^{(1)} \rightsquigarrow v$ is a rainbow path in H_l . \square \square

3.1.4. Reduction: We have gathered all the tools to formulate the reduction. Let the instance for k -subset rainbow connectivity be $\langle G = (V, E), P \rangle$. We construct a graph $\langle G' \rangle$ as an instance for k -rainbow connectivity as follows:

- 1) Construct $H_k = (V_k, E_k)$.
- 2) Let $V' = W_k$, $E' = E_k \cup \{(v_{i,k}, v_{j,k}) : (v_i, v_j) \in E\}$.

$G' = (V', E')$ is the required graph. We call the induced subgraph of G' containing the base vertices as base graph G_k . It can be seen that G_k is isomorphic to G .

Let $(v_{i,k}, v_{j,k}) \in P$. Consider a path $\mathcal{P} = v_{i,k} - x_1 \rightsquigarrow x_s - v_{j,k}$ between $v_{i,k}$ and $v_{j,k}$ of length k or less. All the edges in \mathcal{P} cannot be in H_k from Lemma 3.1. Hence \mathcal{P} contains at least one edge from G_k . Consider a subpath $\mathcal{P}' = x \rightsquigarrow y$ of \mathcal{P} such that x, y are base vertices and all the edges in \mathcal{P}' are in H_k . If $\mathcal{P}' \neq \emptyset$ then $\mathcal{P} \neq \mathcal{P}'$ since \mathcal{P} contains at least one edge from G_k . From Lemma 3.1 and Lemma 3.2, \mathcal{P}' is of length at least k . Hence, length of \mathcal{P} is definitely more than k , contradicting the fact that length of \mathcal{P} is k or less. Hence $\mathcal{P}' = \emptyset$ and all the edges in \mathcal{P} are in G_k . Thus, any path of length k or less between any pair of vertices in P is entirely contained in G_k itself.

We now prove the main theorem of this section which proves that k -rainbow connectivity problem is NP-hard.

Lemma 3.4. *G is k -subset rainbow connected iff G' is k -rainbow connected.*

Proof. “If”: Let $\chi : G' \rightarrow \{c_1, \dots, c_k\}$ be an edge coloring of G' with k colors such that every pair of vertices in G' have a rainbow path between them. We define χ' to be a k -edge coloring of G as follows: $\chi'(v_i, v_j) = \chi((v_{i,k}, v_{j,k}))$ if the edge $(v_{i,k}, v_{j,k}) \in G_k$. By the observations made prior to this lemma, all paths between $v_{i,k}, v_{j,k}$ of length k or less are in G_k itself. Hence the entire rainbow path between $v_{i,k}$ and $v_{j,k}$ must lie in G_k itself. Correspondingly, there is a rainbow path between v_i and v_j in G .

“Only if”: If G is k -subset rainbow connected then let χ be a edge coloring of G such that $\chi : E \rightarrow \{c_1, \dots, c_k\}$ makes it subset rainbow connected. Also, let $\chi' : W_k \rightarrow \{c_1, \dots, c_k\}$ be an edge coloring of H_k such that all pairs of vertices in $(W_k \times W_k) \setminus P_k$ are rainbow connected. We now define an edge coloring $\chi'' : E' \rightarrow \{c_1, \dots, c_k\}$ in G' using χ and χ' as follows. Let $e \in G'$. If $e \in H_k$ then $\chi''(e) = \chi'(e)$ else $\chi''(e) = \chi(e)$. The rainbow paths between every pair of vertices in P_k will be in G_k itself. From Lemma 3.3, all the pairs of vertices in $V' \setminus P_k$ will be rainbow connected. Since there is a rainbow path between every pair of vertices in G' , the graph G' is rainbow connected. \square \square

Therefore, we now conclude the following theorem.

Theorem 3.5. *For every $k \geq 3$, deciding whether $rc(G) \leq 3$ is NP-hard.*

4. CONCLUSION

In this paper, we showed that given any natural number k and a graph G it is NP-hard to determine whether $rc(G) \leq k$. We also show given a bipartite graph, the problem of determining whether $src(G) \leq k$ is NP-hard.

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